

A NOTE ON MINIMAL MATRIX REPRESENTATION OF CLOSURE OPERATIONS

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A matrix M with n columns represents a closure operation $F(A)$, ($A \subseteq X$, $|X|=n$) if for any A , any two rows equal in the columns corresponding to A are also equal in $F(A)$. Let $m(F)$ be the minimum number of rows of the matrices representing F . Lower and upper estimates on $\max m(F)$ are given where \max runs over the set of all closure operations on n elements.

Put $X = \{1, 2, \dots, n\}$. A function $F: 2^X \rightarrow 2^X$ satisfying the conditions

- (1) $A \subseteq F(A)$,
- (2) $A \subseteq B \Rightarrow F(A) \subseteq F(B)$,
- (3) $F(F(A)) = F(A)$, $A, B \subseteq X$

is called a *closure operation*. If M is an $m \times n$ matrix then $F_M(A)$, $A \subseteq X$, contains the i th column of M iff

(4) any two rows identical in columns belonging to A are also equal in the i th column.

It is easy to see that $F_M(A)$ is a closure operation. We say that M represents the closure operation F if $F = F_M$. It is known [1] that any closure operation is representable with an appropriate matrix M . Let $m(F)$ denote the minimum number of rows of these M . Finally, let $R(n) = \max m(F)$ where the maximum is taken over all the closure operations on n elements.

If F is a closure operation define

$$\mathcal{K}_F = \{A: F(A) = X \text{ and } A \text{ is minimal for this property}\}.$$

We say that an $m \times n$ matrix M represents the family \mathcal{K} iff $\mathcal{K} = \mathcal{K}_{F_M}$. Let $m(\mathcal{K})$ denote the minimum number of rows of these M . Finally, let $r(n) = \max m(\mathcal{K})$ where the maximum is taken over all such possible families \mathcal{K} on n elements.

Observe that $m(\mathcal{K}_F) \leq m(F)$ holds for any F . Hence, by definition,

$$(5) \quad r(n) \leq R(n)$$

follows. On the other hand [2] proves the theorem

$$(6) \quad \frac{1}{n^2} \binom{n}{\lfloor n/2 \rfloor} \cong r(n) \cong \binom{n}{\lfloor n/2 \rfloor} + 1.$$

The aim of this paper is to prove the analogous inequalities for $R(n)$:

Theorem.

$$(7) \quad \frac{1}{n^2} \binom{n}{\lfloor n/2 \rfloor} \cong R(n) \cong \left(1 + O\left(\frac{1}{\sqrt{n}}\right)\right) \binom{n}{\lfloor n/2 \rfloor}.$$

Proof. The left hand side follows from (6) and (5). To prove the right hand side we construct a matrix M with this many rows for any given closure operation F .

Let $\mathcal{B}_F = \{A: F(A) = A\}$. It is easy to see [1] that \mathcal{B}_F has the property that $A, B \in \mathcal{B}_F \Rightarrow A \cap B \in \mathcal{B}_F$ and conversely if \mathcal{B} possesses this property then there is an F with $\mathcal{B}_F = \mathcal{B}$ (F is defined by $F(A) = \bigcap B$ where the intersection is taken over all sets B satisfying $A \subseteq B \in \mathcal{B}$). If \mathcal{A} is any family of sets let $\tilde{\mathcal{A}}$ denote the family of all possible intersections of its members. Given F , we form a sequence $\mathcal{B}_F = \mathcal{B}_0 \supseteq \mathcal{B}_1 \supseteq \mathcal{B}_2 \supseteq \dots \supseteq \mathcal{B}_k$ of families. Choose three different members (if there exist) A, B and C of \mathcal{B}_i satisfying $A \cap B = C$. \mathcal{B}_{i+1} is defined by $\mathcal{B}_{i+1} = \mathcal{B}_i - \{C\}$. The statement $\tilde{\mathcal{B}}_{i+1} = \tilde{\mathcal{B}}_i$ is trivial. Therefore, $\tilde{\mathcal{B}}_i = \tilde{\mathcal{B}}_0 = \mathcal{B}_F$ holds for any i ($0 \leq i \leq k$). It is easy to see that the procedure stops after a finite number of steps: in \mathcal{B}_k there are no 3 different members with $A \cap B = C$. A theorem of Kleitman [3] states that the size of a family of subsets of an n element set containing no 3 different members with $A \cap B = C$ cannot exceed $\left(1 + O\left(\frac{1}{\sqrt{n}}\right)\right) \binom{n}{\lfloor n/2 \rfloor}$.

We construct M using the family $\mathcal{B}_k = \{A_1, \dots, A_s\}$. The j th ($1 \leq j \leq n$) entry of the i th ($0 \leq i \leq s$) row is

$$a_{ij} = \begin{cases} 0 & \text{if } i = 0 \\ 0 & \text{if } 1 \leq i \leq s, \quad j \in A_i \\ i & \text{if } 1 \leq i \leq s, \quad j \notin A_i. \end{cases}$$

The number of rows of M is $1 + |\mathcal{B}_k| \cong \left(1 + O\left(\frac{1}{\sqrt{n}}\right)\right) \binom{n}{\lfloor n/2 \rfloor}$. We have to prove only that M represents F . Let A be an arbitrary subset of X . Suppose first that $j \in F(A)$. Two rows i_1, i_2 can be equal in the columns corresponding to A only then if all these entries are zeros. That is, if $A \subseteq A_{i_1}, A_{i_2}$, (2) implies $F(A) \subseteq F(A_{i_1}), F(A_{i_2})$. As $A_{i_1}, A_{i_2} \in \mathcal{B}_F$, $F(A) \subseteq A_{i_1}, A_{i_2}$ follows. Consequently, the two rows are equal ($0=0$) in the j th column.

Suppose now $j \notin F(A)$, $F(A) \in \mathcal{B}_F$ holds by (3). $\tilde{\mathcal{B}}_k = \mathcal{B}_F$ implies that $F(A)$ is an intersection of some (≥ 1) members of \mathcal{B}_k . One of these members, say A_i does not contain j . The 0th and the i th rows are equal in A but are different in the j th column. That is, $F_M = F$ holds. ■

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References

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